

Constraint programming based iterative heuristic for scheduling trains and maintenance tasks

Team 2.7.2 solution

S. Carpov L. Cudennec F. Galea

CEA, LIST

2016 RAS problem solving competition

Our team

- ▶ Researchers at CEA (LIST Institute)
- ▶ CEA – French Alternative Energies and Atomic Energy Commission
 - ▶ 10 research centers across France
 - ▶ >15.000 people
 - ▶ Research areas:
 - ▶ nuclear energy (fission and fusion)
 - ▶ **technological research for industry**
 - ▶ defense and security
 - ▶ fundamental research in physical and life sciences.

Outline

RAS problem solving competition

Our Solution

Experimentation and results

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Routing Trains through a Railway Network

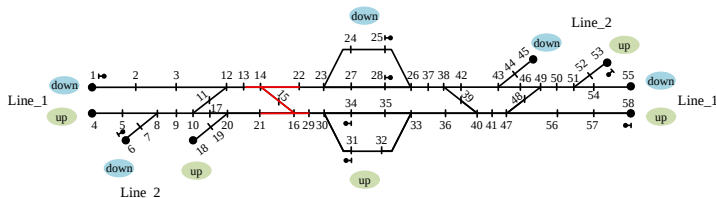
Joint optimization on train timetabling and maintenance task scheduling

► Problem input:

- railway network description
 - nodes, links, cells, blocks, stations
- trains to route and timetable
 - origin, destination, time window, stops
- maintenance tasks
 - duration, cells to maintain, time window

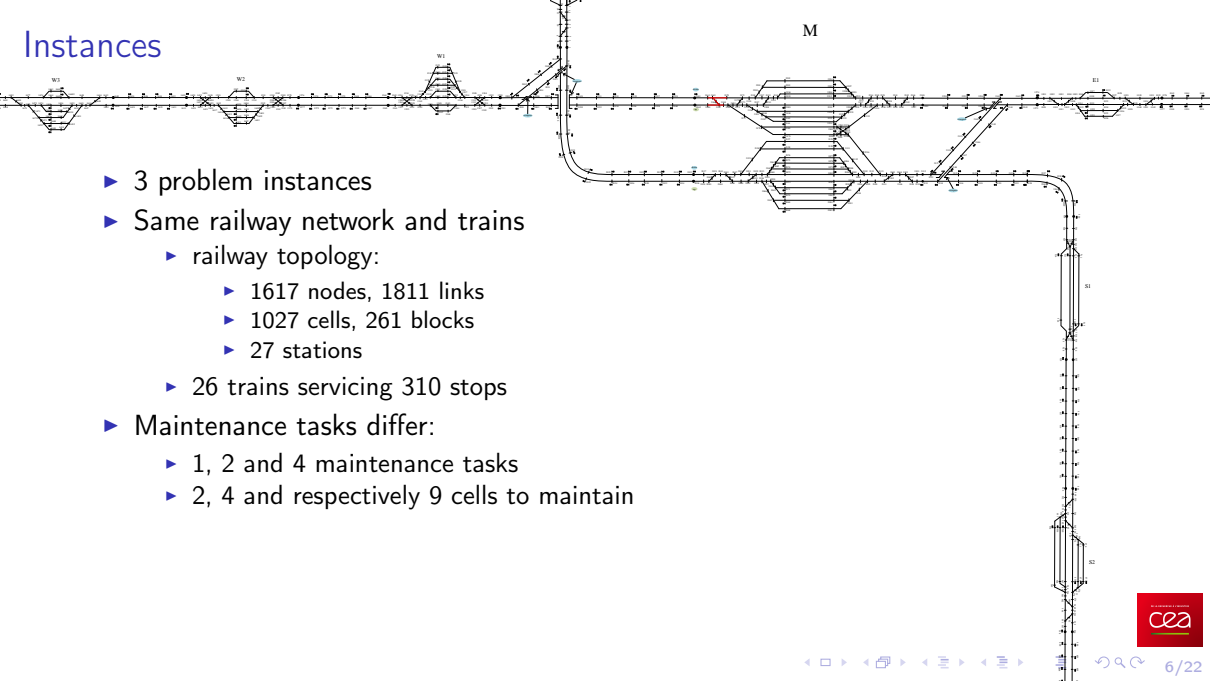
► Goal:

- train timetabling
 - railway network routes
 - node arrival times
 - maintenance task scheduling
- Objective: minimize total travel time of all trains



Instances

- ▶ 3 problem instances
- ▶ Same railway network and trains
 - ▶ railway topology:
 - ▶ 1617 nodes, 1811 links
 - ▶ 1027 cells, 261 blocks
 - ▶ 27 stations
 - ▶ 26 trains servicing 310 stops
- ▶ Maintenance tasks differ:
 - ▶ 1, 2 and 4 maintenance tasks
 - ▶ 2, 4 and respectively 9 cells to maintain



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Railway network graph

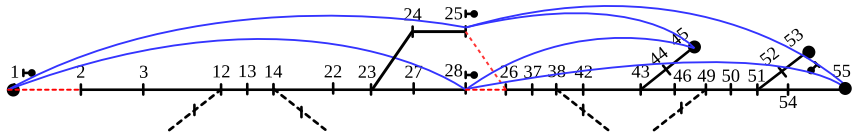
Definition

- ▶ A graph $G = (N, L, W)$:
 - ▶ N set of railway nodes
 - ▶ L set of railway links and *block traversal links* (introduced on next slide)
 - ▶ W set of link weighting functions
 - ▶ $W = \{c_t : t \in \mathcal{T}\}$ and $c_t : L \rightarrow \mathbb{N}$
- ▶ Weighting function $c_t(l)$ gives link $l \in L$ traversal time:
 - ▶ depends on train speed multiplier and link speed limit
 - ▶ includes minimum dwelling time for station siding tracks
 - ▶ infinite value for station main tracks where train t should stop

Railway network graph

Block traversal links

- ▶ Railway network block sections give authorized ways to traverse the network
- ▶ Block traversal links:
 - ▶ artificial links corresponding to the traversal of blocks
 - ▶ block traversal link weight is the sum of weights of corresponding links
 - ▶ block first link removed from G



Railway network graph

Paths

- ▶ Valid railway network traversal routes with respect to constraints:
 - ▶ link minimum running time
 - ▶ train required stops
 - ▶ minimum dwelling time
 - ▶ block section selection
- ▶ Correspond to train routes with no maintenance tasks and independent train traversals
 - ▶ i.e. G path lengths are lower bounds to actual train travel times

Global lower bound

- ▶ Global lower bound on total travel times for all the trains:

$$GLB = \sum_{t \in T} short_path_len(G, orig(t), dest(t), c_t)$$

- ▶ $short_path_len(G, s, e, c_t)$ – shortest path length from node s to node e using weighting function c_t
- ▶ Gives total travel time for a solution not considering train collision and maintenance tasks
- ▶ We are looking for solution values as close as possible to GLB

Simple Heuristic

General idea

Decompose problem solving in two steps:

1. Route trains through the railway network
2. Schedule trains and maintenance tasks using a constraint programming (CP) model

Heuristic first-step

Railway network train routing

1. Build an auxiliary railway network graph G' (built from G)
 - ▶ Restrict “pessimistically” speed limits of links impacted by maintenance tasks
 - ▶ I.e. use the worst case link traversal time
 - ▶ Simplified problem model where all trains always traverse maintained links at low speed
 - ▶ Restrict middle station traversal by train origin and destination
 - ▶ Reduces congestion and balances train traversal of middle station
2. Find for each train the shortest path in graph G'

CP model

Interval variables (IV)

Interval variable definition

- ▶ decision variable whose value is an interval of integers
- ▶ modeled using 2 values: start and end of the interval

Employed IVs

- ▶ Train link traversal
 - ▶ path links found in the first step
- ▶ Train block¹ traversal
 - ▶ block interval variables “span” over respective link interval variables
- ▶ Maintenance task
 - ▶ maintained cells are grouped => one interval variable per maintenance task

¹To ease the exposition all cells are considered blocks.

CP model

Constraints and objective

- ▶ Link traversal and maintenance tasks minimal duration
- ▶ Order of link traversal (respect train paths)
- ▶ Starting time windows of trains and maintenance tasks
- ▶ No overlap
 - ▶ Same link for different trains IVs
 - ▶ Same block for different trains IVs ²
 - ▶ Block and maintenance tasks IVs with common cells
- ▶ Objective: minimize the sum of link IV durations

²Special case for arrival blocks.

Iterative heuristic

- ▶ Second step of simple heuristic is a cumbersome process
 - ▶ Large CP model for off-the-shelf solvers
 - ▶ Poor search space exploration
- ▶ Main idea of iterative heuristic:
 - ▶ fix the schedule for some trains and solve the reduced CP model
 - ▶ do this iteratively
- ▶ Instead of solving a large CP problem solve several easier sub-problems

Iterative heuristic

Iterative heuristic steps

1. Find a starting solution (simple heuristic, small time limit)
2. Repeat several times:
 - 2.1 Choose a set of trains to schedule
 - 2.2 Fix other trains schedule and solve the reduced CP model

Step 2.1 – train to schedule choice

- ▶ Use gap between train travel time and its lower bound to make decision:
 - A choose trains having the largest gap
 - B privilege trains with common traversal segments besides lower bound gap

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Experimentation

- ▶ python for heuristics first-step and *GLB*
- ▶ ILOG CP solver for constraint programming models
- ▶ Execute on a single core of an AMD Opteron 6172 processor (2.1GHz)
- ▶ Simple and iterative heuristics time limit – 2 hours
- ▶ *GLB* used to assert solution quality
- ▶ Beer & pizzas

Results

Total travel time of all trains

Heuristic	Solution value in seconds (gap to <i>GLB</i>)		
	Case 1 1 MOT	Case 2 2 MOTs	Case 3 4 MOTs
Simple	157459 (1.83%)	158339 (2.40%)	161341 (4.34%)
Iterative A	156096 (0.95%)	156865 (1.45%)	160608 (3.87%)
Iterative B	155294 (0.43%)	157176 (1.65%)	158985 (2.82%)
<i>GLB</i>	154625		

- ▶ *GLB* proves to be rather powerful on given instances
- ▶ Less maintenance tasks => smaller solution gap to *GLB*

CP model generalization ideas

- ▶ Consider several traversal ways of stations (e.g. stations with >1 siding track)
 - ▶ allows to switch heading train
- ▶ Several possible paths for train traversal (generalization of previous case)
 - ▶ e.g. dynamic choice of middle station traversal path
- ▶ Use one IV for each cell of maintenance tasks (do not group maintained cells)
 - ▶ adds flexibility to the model

Thanks!